

Even when we don't have an explicit formula for a recurrence, we can often work out valuable information about it using inductive methods:

**Proposition 1.** *Let  $b_n$  be given by a recurrence relation:  $b_1 = 1$ ,  $b_2 = 2$ , and  $b_n = nb_{n-1} + b_{n-2}$  for  $n > 2$ . For positive integers  $n$ ,  $2 \mid b_n$  if and only if  $2 \mid n$ .*

*Proof.* Let  $P(n)$  be the assertion that  $2 \mid b_n$  if and only if  $2 \mid n$ ; alternatively, it could be phrased as the assertion that  $b_n$  and  $n$  have the same parity. The base cases  $P(1)$  and  $P(2)$  are easily verified:  $P(1)$  is odd as is 1, and  $P(2)$  is even as is 2.

Now let us fix a  $k \geq 2$ , and inductively assume that  $P(1), P(2), \dots, P(k)$  are all true. Specifically, we will use the facts that  $P(k)$  and  $P(k-1)$  are true below. Now we have two cases: either  $2 \mid (k+1)$ , in which case our goal is to show that  $2 \mid b_{k+1}$ , or  $2 \nmid (k+1)$ , in which case our goal is to show that  $2 \nmid b_{k+1}$ . We address these casewise:

**Case I:  $2 \mid k+1$  (a.k.a.  $k+1$  is even).** Since  $k+1$  is even,  $k$  is odd and  $k-1$  is even, and since we presumed  $P(k)$  and  $P(k-1)$  to be true, it thus respectively follows that  $b_k$  is odd and  $b_{k-1}$  is even. Then, we note that  $b_{k+1} = (k+1)b_k + b_{k-1}$ ; this is a sum of a product of an even and odd number with an even number; the result of such a computation is even, so  $2 \mid b_{k+1}$ .

**Case II:  $2 \nmid k+1$  (a.k.a.  $k+1$  is odd).** Since  $k+1$  is odd,  $k$  is even and  $k-1$  is odd, and since we presumed  $P(k)$  and  $P(k-1)$  to be true, it thus respectively follows that  $b_k$  is even and  $b_{k-1}$  is odd. Then, we note that  $b_{k+1} = (k+1)b_k + b_{k-1}$ ; this is a sum of a product of an odd and even number with an odd number; the result of such a computation is odd, so  $2 \nmid b_{k+1}$ .  $\square$

We can try this with an even more complicated recurrence. Here's one which doesn't have a specific order, but *is* a recurrence relation, and is actually a quite significant set of numbers:

**Definition 1.** The *Catalan numbers*  $C_1, C_2, C_3, \dots$  are given by the recurrence relation:

$$C_1 = 1$$

$$C_n = \sum_{i=1}^{n-1} C_i C_{n-i} \text{ for } n > 1$$

so, for example, we can work out the first few Catalan numbers:

$$C_1 = 1$$

$$C_2 = C_1 C_1 = 1 \cdot 1 = 1$$

$$C_3 = C_2 C_1 + C_1 C_2 = 1 \cdot 1 + 1 \cdot 1 = 2$$

$$C_4 = C_3 C_1 + C_2 C_2 + C_1 C_3 = 2 \cdot 1 + 1 \cdot 1 + 1 \cdot 2 = 5$$

$$C_6 = C_4 C_1 + C_3 C_2 + C_2 C_3 + C_1 C_4 = 5 \cdot 1 + 2 \cdot 1 + 1 \cdot 2 + 1 \cdot 5 = 14$$

and the next few are 42, 132, 429, 1430.

We can prove a couple of things about the Catalan numbers. For instance, their growth rate appears to be quite fast, which we can quantify pretty painlessly by showing that it's at least exponential:

**Proposition 2.** *For every positive integer  $n$ ,  $C_n \geq 2^{n-2}$ .*

*Proof.* We shall separately demonstrate the truth of this statement for  $n = 1$ , since  $C_1 = 1 \geq 2^{-1}$ ; henceforth we proceed by induction with the base case  $n = 2$ . We prove the base case arithmetically:  $C_2 = 1 = 2^0$ , so this inequality holds when  $n = 2$ . We now proceed to the inductive step.

We fix a  $k$ , and may then accept the hypothesis that  $C_i \geq 2^{i-2}$  for all  $i \leq k$ . We now want to place a lower bound on the value of  $C_{k+1}$ , which we can do since  $C_{k+1}$  is calculated based on the values of various smaller  $C_i$ .

$$\begin{aligned} C_{k+1} &= C_1 C_k + C_2 C_{k+1} + \cdots + C_k C_1 \\ &\geq C_1 C_k + C_k C_1 \\ &\geq 1 \cdot C_k + C_k \cdot 1 \\ &\geq 2C_k \geq 2(2^{k-2}) = 2^{(k+1)-2} \end{aligned}$$

□

In actuality even better upper bounds on the Catalan numbers are possible.

Playing with these further, we might discover another neat property of Catalan numbers:

**Conjecture 1.** *For a positive integer  $n$ ,  $C_n$  is odd if and only if  $n$  is a power of 2.*

Proving this is actually quite easy, with the use of a helpful lemma.

**Lemma 1.** *For every positive integer  $n$ ,  $C_{2n}$  has the same parity as  $C_n$ .*

and then we can prove our conjecture easily by use of this fact.

**Proposition 3.** *For every positive integer  $n$ ,  $C_n$  is even if and only if  $n$  is a power of 2.*

*Proof.* For clarity, let us restate exactly what we want to prove for each positive integer  $n$ : we wish to show that if  $n$  is a power of 2, then  $C_n$  is odd, and otherwise,  $C_n$  is even. We shall prove this by induction on  $n$ . The cases  $n = 1$  and  $n = 2$  are obviously true since  $C_1$  and  $C_2$  are both odd (and 1 and 2 are powers of 2).

Fixing  $k \geq 2$ , we may assume the inductive hypothesis that the parities of each  $C_i$  for  $1 \leq i \leq k$  follows the observed pattern. Now, we must divide into cases:

**Case I:  $k + 1$  is a power of 2.** Since  $k + 1 > 2$  and  $k + 1 = 2^r$  for some positive integer  $r$ , it is easy to see that  $r > 1$  and thus  $k + 1 = 2(2^{r-1})$ . By our lemma,  $C_{k+1}$  and  $C_{2^{r-1}}$  have the same parity, and by our inductive hypothesis, since  $2^{r-1}$  is a power of 2,  $C_{2^{r-1}}$  is odd.

**Case II:  $k + 1$  is an even number which is not a power of 2.** Since  $k + 1$  is even and positive,  $k + 1 = 2\ell$  for some positive integer  $\ell$ . By our lemma,  $C_{k+1}$  and  $C_\ell$  have the same parity. Since  $k + 1$  is not a power of 2,  $\ell$  is not a power of 2, since if it were the case that  $\ell$  were a power of 2, so would  $k + 1$  be. Thus, by our inductive hypothesis,  $C_\ell$  is even, so  $C_{k+1}$  is even.

**Case III:  $k + 1$  is an odd number greater than 2 (and thus not a power of 2).** Then  $k + 1 = 2\ell + 1$  for some integer  $\ell$ , and the recurrence for  $C_{k+1}$  is:

$$C_{k+1} = C_1 C_k + C_2 C_{k-1} + \cdots + C_\ell C_{\ell+1} + C_{\ell+1} C_\ell + \cdots + C_{k-1} C_2 + C_k C_1$$

which can be rearranged into  $C_{k+1} = 2(C_1 C_k + C_2 C_{k-1} + \cdots + C_\ell C_{\ell+1})$ , so  $C_{k+1}$  is even. □

## 1 Induction and finite sets

One can prove things by induction on finite sets. For instance, take the problem on the first pset.

Give that old saw about how all horses are the same color, unless there's a volunteer.